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Enabling GPU support for the COMPSs-Mobile framework

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4th Workshop on Accelerator
Programming Using Directives

COMPSS-Mobile infrastructure



Programming Model: COMPSs

- Programming
 - Sequential code
 - Standard programming languages with no APIs
 - Infrastructure agnostic
 - Task-based
 - An annotated interface contains the declaration of those methods (CE) selected to become tasks with information about the accesses to data
- Execution managed by a runtime system
 - Replaces CE invocations by asynchronous tasks
 - Monitors data-dependencies among tasks
 - Orchestrates the execution on the underlying infrastructure
 - Transfers data values
 - Submits task executions

Programming Model Example

Matmul.java

```
public class Matmul {  
    public static void main(String[] args) {  
        int[][] A;  
        int[][] B  
        int[][] C;  
        ...  
        C = multiply(A, B);  
        ...  
    }  
    public static int[][] multiply (int[][] A, int[][] B) {  
        // Matrix multiplication code  
        //C = AB  
        ...  
        return C;  
    }  
}
```

CEI.java

```
public interface CEI {  
    @Method(declaringClass="es.bsc.compss.matmul.Matmul")  
    int[][] multiply (  
        @Parameter(direction = IN)  
        int[][] A,  
        @Parameter(direction = IN)  
        int[][] B  
    );  
}
```

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Programming Model Extension

- Add the OpenCL kernel codes in application resource files
- Indicate on the CEI the CE implementation as an OpenCL kernel using the **@OpenCL** annotation indicating:
 - Resource containing the code implementation (kernel)
 - Number of work-items (globalWorkSize)
 - Work-items ID offset (offset)
 - Work groups size (localWorkSize)
 - Result Size (result)

Programming Model Extension Example

Matmul.java

```
public class Matmul {  
    public static void main(String[] args) {  
        int[][] A;  
        int[][] B;  
        int[][] C;  
        ...  
        C = multiply(A, B);  
        ...  
    }  
    public static int[][] multiply (int[][] A, int[][] B) {  
        // Matrix multiplication code  
        //C = AB  
        ...  
        return C;  
    }  
}
```

Matmul.cl

```
__kernel void multiply (  
    __global const int *a,  
    __global const int *b,  
    __global int *c) {  
    //Matrix multiplication code  
    // C = AB  
    ...  
}
```

CEI.java

```
public interface CEI {  
    @OpenCL(kernel="matmul.cl", globalWorkSize="par0.x,par1.y", resultSize="par0.x,par1.y")  
    @Method(declaringClass="es.bsc.compss.matmul.Matmul")  
    int[][] multiply (  
        @Parameter(direction = IN)  
        int[][] A,  
        @Parameter(direction = IN)  
        int[][] B  
    );  
}
```

Kernel Execution Lifecycle

[0 .- Fetch input values from a remote node]

1 .- Create a cl_buffer for each parameter and the result

```
cl_mem par0Dev = clCreateBuffer(...);
```

2 .- Fill the cl_buffers corresponding to the input values

```
clEnqueueWriteBuffer(..., par0Dev, ...);
```

3 .- Run the kernel

```
cl_kernel task1 = clCreateKernel (...);  
clSetKernelArgument(task1, 0, par0Dev);  
clNDRangeKernel(..., task1, ...)
```

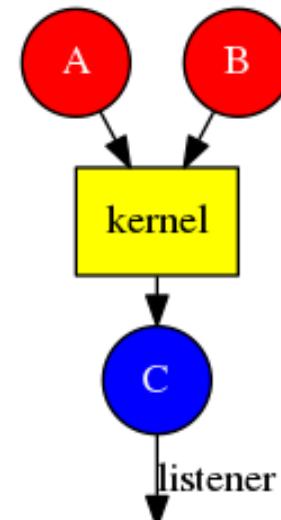
4 .- Bring back the modified/result values

```
clEnqueueReadBuffer(..., par0Dev, ...)
```

5 .- Notify the existence of the created values to enable data-dependent tasks executions.

Kernel Execution Lifecycle Management

- Management is twofold:
 - The runtime library directly controls the creation of OpenCL structures and Host memory management
 - Fetching data from remote workers
 - Allocating the necessary device memory space to host the buffers
 - Creating the Kernel object and setting up its arguments
 - Publishing the value existence to release data dependencies
 - It delegates on the out-of-order mode of OpenCL the execution of OpenCL commands:
 - `clEnqueueWriteBuffer`
 - `clNDRangeKernelExecution`
 - `clEnqueueReadBuffer`

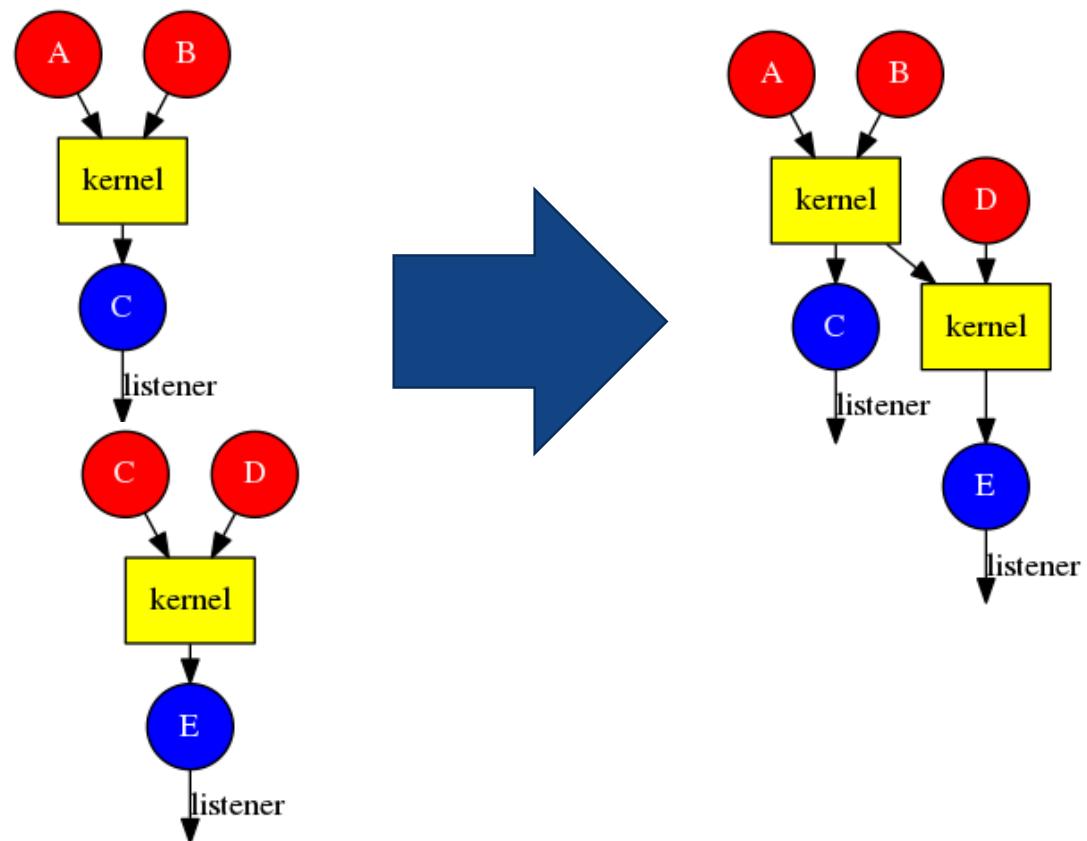


Optimizations

- Keep values on the device memory to reuse them
- Keep track of the events generating the values

$C = \text{matmul}(A, B)$

$E = \text{matmul}(C, D)$



Executing Platform Selection

- Each task is profiled to measure the execution time and energy consumption
- Using historical values of previous executions the runtime forecasts:
 - The execution time on the resource
 - The energy consumption of the execution
 - The waiting time before the execution start
- According to a heuristic defined by the application user and based on the forecasts, the runtime picks one platform

Performance Evaluation

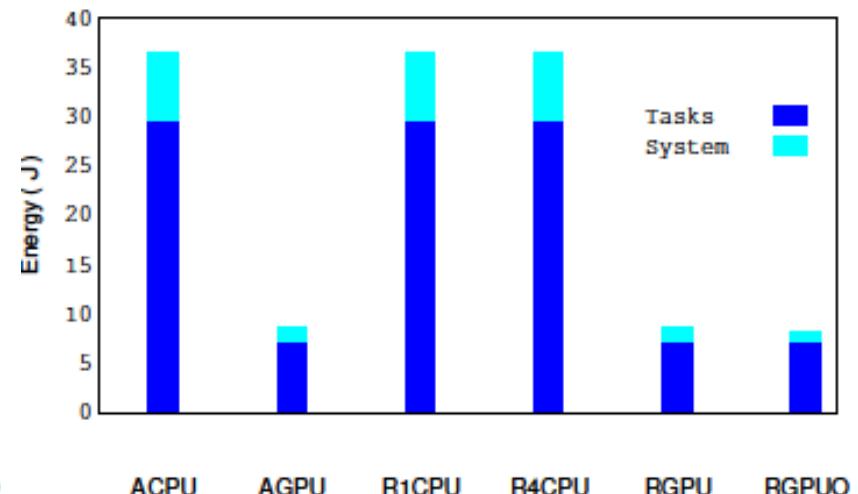
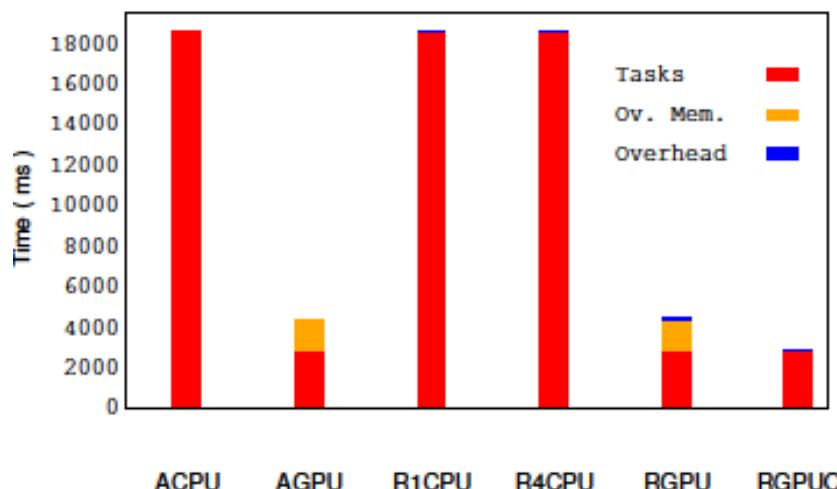
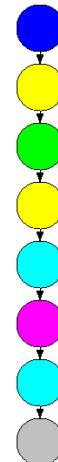


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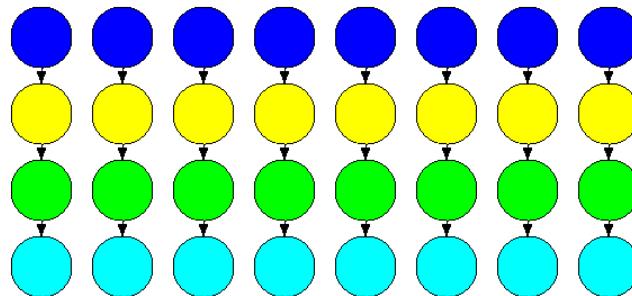
Digits Recognition

- CNN recognizing 512 handwritten digits
- Sequence of 8 steps (tasks) processing all 512 images
- GPU speeds up task calculation by 6.5x
- Optimizations reduce the 1,531 ms overhead to 5 ms
- Energy consumption falls from 30 J to 8J



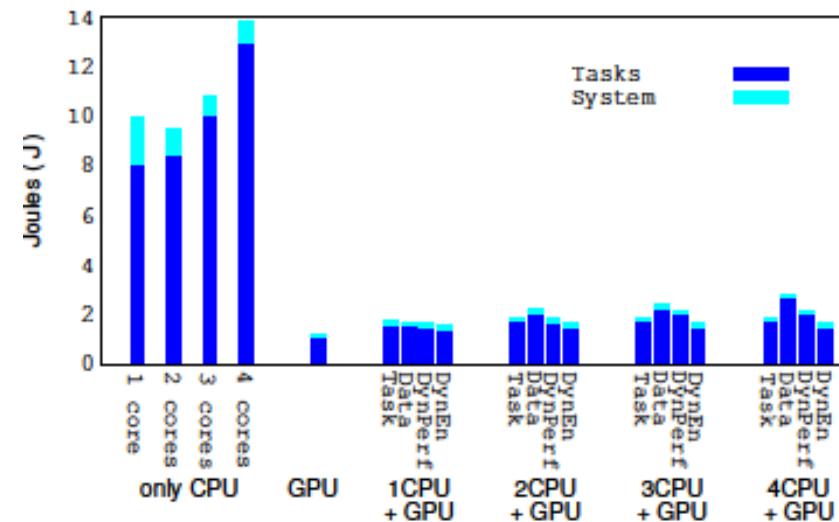
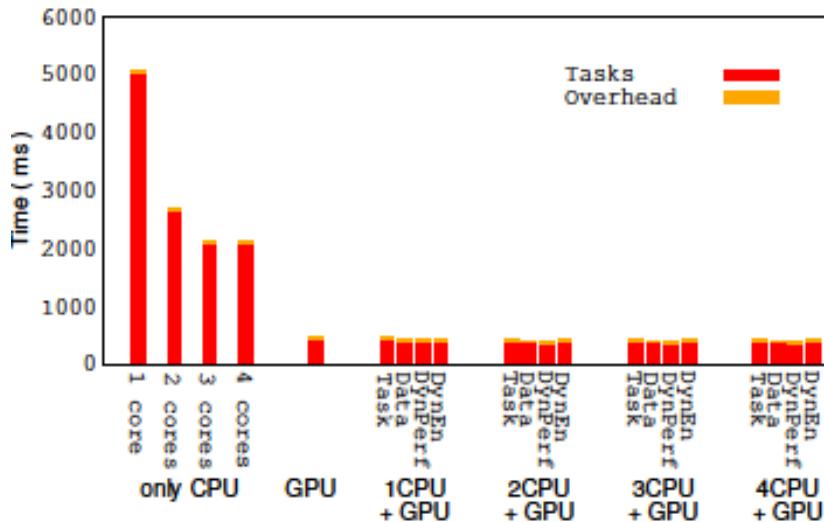
Canny Edge detection

- Detects the edges in 30 frames of a video
- Each frame goes through 4 stages (4 tasks)



- Developers normally try to balance the load according to:
 - Data Partitioning: assign the processing of a whole frame to a resource
 - Task Partitioning: run the two first tasks of every frame on the GPU and the last two on the CPU
- 2 dynamic policies
 - DynPerf: minimizes the execution time
 - DynEn: slows down 1 ms the execution if that saves 5 mJ

Canny Edge detection



- GPU speeds up 12x the task computation
- GPU reduces the energy consumption from 9.89 J to 1.34 J
- Data Partitioning achieves lower execution times but consumes more energy than Task Partitioning
- DynPerf achieves the lowest execution time (375 ms)
- DynEn gets lower energy consumptions while its performance compares to Task Partitioning

Conclusions

- The programming model allows developers to code parallel applications that run on heterogeneous systems being unaware of the parallelization and collaboration issues
- Using GPUs reduces drastically the timespan and the energy consumption of the application
- Dynamic policies provide applications with
 - Portability
 - Flexibility
- Code is available at
 - <http://compssdev.bsc.es/gitlab/flordan/WACCPD>
 - <http://compssdev.bsc.es/gitlab/flordan/COMPSS-Mobile>

Future Work

- Enable the usage of the GPUs to compute tasks on remote nodes
- Improve the forecasts for the energy and end time
- Allow the dynamic policies to change an initial decision on which resources the task will run



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Thank you!

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